



# Hiring Janitors

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# Introduction

- A museum has open job spots for janitors and is currently hiring. They would like to know:
  - How many janitors they need to hire.
  - A work schedule for the janitors.
  - An algorithm to find both using the number of guests.

# Problem Statement

- Everyday,  $x$  guests come to the museum, where  $1000 \leq x \leq 1500$ .
- Guests produce  $99x/100$  total trash per day.
- One janitor can clean 100 trash in 1 hour; the maximum work day is 4 hours per day, and the maximum work week is 6 days.
- The max number of janitors is 25, and the minimum number is 1.
- Find an algorithm using the number of people and the work hours to find the proper number of janitors required to clean the trash.

# Results p1

- First results:

- The first equation I found was  $j = \left(\frac{\frac{100x}{100}}{th}\right)$

- List of variables:

- h=hours a day
- j=#of janitors.
- x=#of guests.
- a=hrs on the job so far.
- y=total janitors.
- l is janitors off.
- t=#of trash cleanable in 1 hour(important for fatigue),t=100(for now).
- z=extra people coming in every week(important for ad campaign)

VARIABLES:

h=hours a day.                      j=#of janitors.  
x=#of guests.    a=hrs on the job so far.  
y=total janitors.                      l is janitors off.  
t=#of trash cleanable in 1 hour.  
z=extra people coming in every week.  
a=hours previously worked that day.

# Results p2

- First results (cont.):

- Schedule Table to the right:
- The general algorithm for the table:
  - hire at minimum twice the number of janitors for day 1
  - for the next day take  $j$  janitors from off duty to on duty
  - the janitors on duty the previous day get the day off
- In the table, you have janitors 1, 2, 3, and 4 working, and 5 thru 11 resting. For day two, you have 5, 6, 7, and 8 on the job, and 9, 10, 11, 1, 2, 3, and 4 resting.

	y=11											
	Work(j)				off(l)							
s	1	2	3	4		5	6	7	8	9	10	11
m	5	6	7	8		9	10	11	1	2	3	4
t	9	10	11	1		2	3	4	5	6	7	8
w	2	3	4	5		6	7	8	9	10	11	1
t	6	7	8	9		10	11	1	2	3	4	5
f	10	11	1	2		3	4	5	6	7	8	9
s	3	4	5	6		7	8	9	10	11	1	2

VARIABLES:

h=hours a day.                    j=#of janitors.  
 x=#of guests.    a=hrs on the job so far.  
 y=total janitors.                l is janitors off.  
 t=#of trash cleanable in 1 hour.  
 z=extra people coming in every week.  
 a=hours previously worked that day.

# Results p3

- Ad Campaign:
  - The museum starts an ad campaign that brings in  $Z$  more people a week.
  - Equation:  $X_n = Z + X_{(n-1)}$
  - plug into original equation.
  - Make Table.

## VARIABLES:

$h$ =hours a day.                       $j$ =#of janitors.  
 $x$ =#of guests.     $a$ =hrs on the job so far.  
 $y$ =total janitors.                       $l$  is janitors off.  
 $t$ =#of trash cleanable in 1 hour.  
 $z$ =extra people coming in every week.  
 $a$ =hours previously worked that day.



# Results p4

- Fatigue:

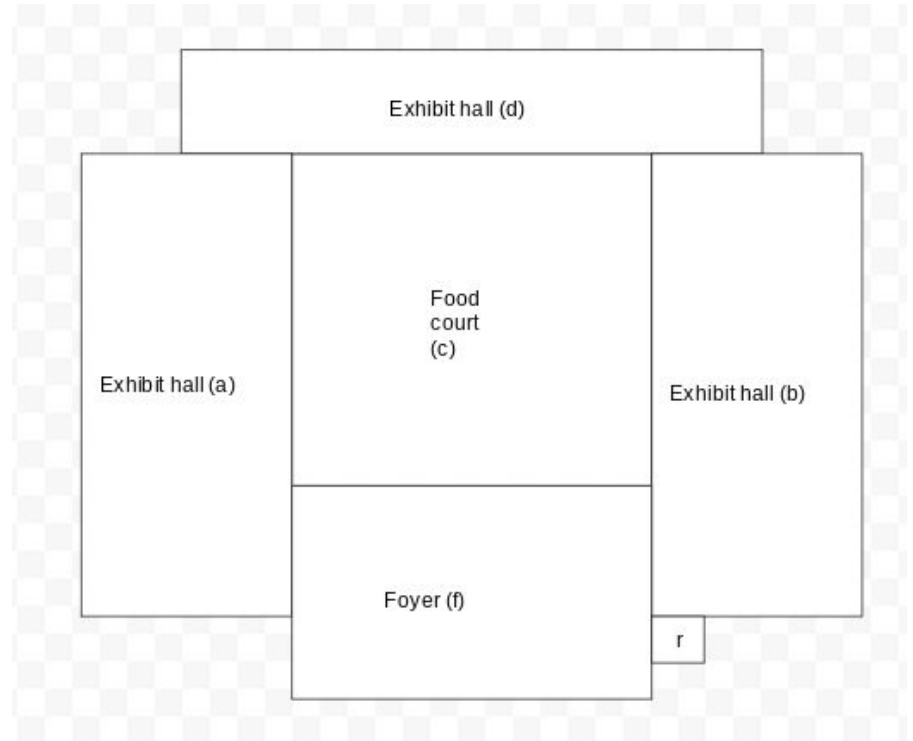
- Fatigue Equation:  $t=100-10a$ 
  - $t$ =cleanable trash per hour for 1 janitor
  - $a$ =number of hours previously worked that day
- For hour 1:
  - $a=0, t=100$
- For hour 2:
  - $a=1, t=90$
- And so on.
- Plug into original equation and then into table algorithm.

VARIABLES:

$h$ =hours a day.                       $j$ =#of janitors.  
 $x$ =#of guests.    $a$ =hrs on the job so far.  
 $y$ =total janitors.                       $l$  is janitors off.  
 $t$ =#of trash cleanable in 1 hour.  
 $z$ =extra people coming in every week.  
 $a$ =hours previously worked that day.

# Results p5

- Path Algorithm p1:
  - The museum wants us to find a path for janitors so that all trash is cleaned and all previously mentioned constraints are followed.
  - This is where we lose a single definitive algorithm.



VARIABLES:  
h=hours a day. j=#of janitors.  
x=#of guests. a=hrs on the job so far.  
y=total janitors. l is janitors off.  
t=#of trash cleanable in 1 hour.  
z=extra people coming in every week.  
a=hours previously worked that day.

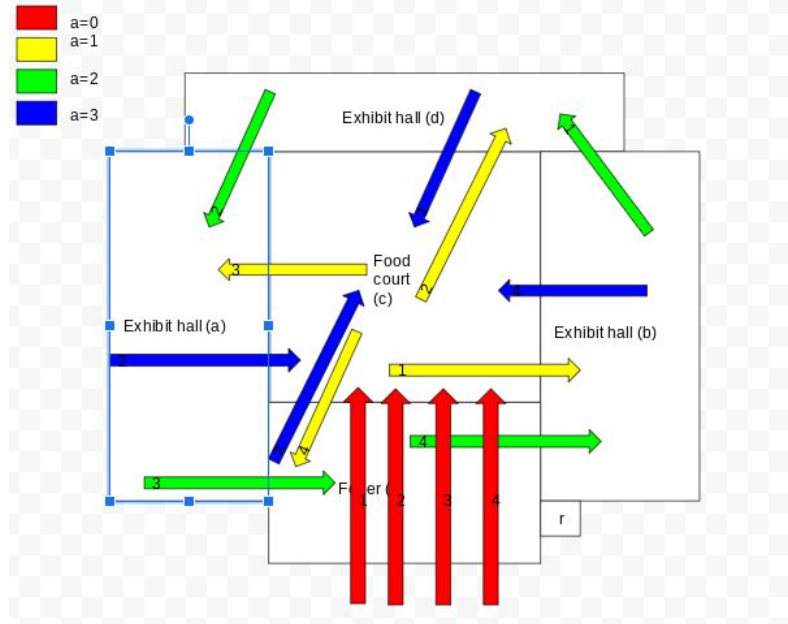


# Results p6

## ● Path Algorithm p2

- Best Algorithm for  $j=4$ ,  $h=4$ , and  $x=1000$ :

- $a=0$ : Send everyone to c.
- $a=1$ : one to a, one to b, one to d, one to f.
- $a=2$ : one to a, one to b, one to d, one to f.
- $a=3$ : all to c.
- When done go to r.



### VARIABLES:

$h$ =hours a day.  $j$ =#of janitors.  
 $x$ =#of guests.  $a$ =hrs on the job so far.  
 $y$ =total janitors.  $l$  is janitors off.  
 $t$ =#of trash cleanable in 1 hour.  
 $z$ =extra people coming in every week.  
 $a$ =hours previously worked that day.

# Results p7

- Path Algorithm p3:

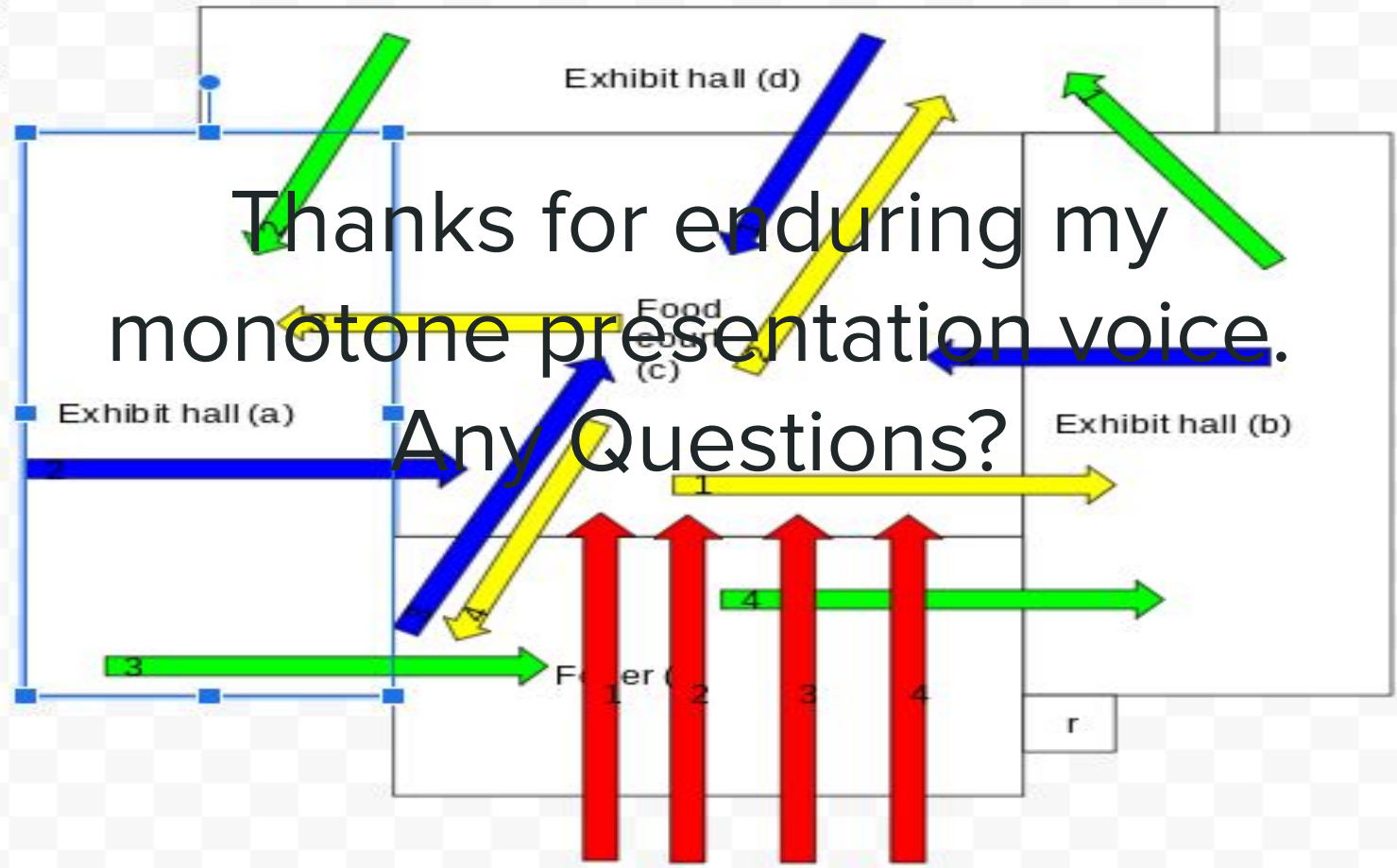
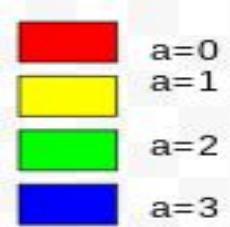
- This general process is applicable for all  $1 \leq y \leq 25$ :
  - Send all janitors to dirtiest room.
  - Send as close to an equal number of janitors to each room as you can get.
  - If any trash left in dirtiest room and sending all janitors to dirty room during last hour does not work, send enough janitors to dirty room as needed during step 2.
  - As stated, and if needed, send all janitors to dirtiest room.
  - If finished go to breakroom.

VARIABLES:

h=hours a day.                      j=#of janitors.  
x=#of guests.    a=hrs on the job so far.  
y=total janitors.                      l is janitors off.  
t=#of trash cleanable in 1 hour.  
z=extra people coming in every week.  
a=hours previously worked that day.

# Conclusion

- It is solvable. Everything up to the path algorithm was very straightforward, but finding paths took much longer than expected.
- What I would do next time:
  - I would probably add in the restriction that janitors can only walk into rooms from adjacent rooms.



Thanks for enduring my monotone presentation voice.

Any Questions?